A stochastic block model for hypergraphs

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SBM for Hypergraphs

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Outline

1 The need for higher-order interactions

2 Stochastic blockmodel for hypergraphs

3 Experiments

4 Conclusions and perspectives

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Higher-order interactions I

Motivations

- Networks or graphs focus on pairwise interactions
- These type of pairwise interactions can already be quite elaborate: undirected/directed, binary/weighted, simple/multiple, static/dynamic, multiplex or multi-layers, ...
- Nonetheless pairwise interactions are not sufficient to describe the nature of complex interactions :
 - e.g. the presence of a 3rd chemical component may modify the interaction of 2 other ;
- Collective interactions or group interactions are richer than just pairwise interactions

 \hookrightarrow These are called higher-order interactions (HOI).

Higher-order interactions II

Where do we find HOI?

- Social networks: triadic and larger groups (as early as Simmel, 1950)
- Scientific co-authorship,
- Interactions between chemical components,
- Interactions between neurons in brain networks,
- etc

These interactions **CAN NOT** be represented by a graph.

Higher-order interactions III

This is a nice recent review (2020):



Networks beyond pairwise interactions: Structure and dynamics

Federico Battiston ^{a,*}, Giulia Cencetti ^b, Iacopo Iacopini ^{c,d}, Vito Latora ^{c,e,f,g}, Maxime Lucas ^{h,i,j}, Alice Patania ^k, Jean-Gabriel Young¹, Giovanni Petri ^{m,n}

Pairwise vs HOI

HOI are defined as sets of interacting entities. e.g. $V = \{a, b, c, d, e\}; \mathcal{I} = \{\{a, b, c\}, \{a, d\}, \{c, d\}, \{c, e\}\}$



Naïve Graph representation: clique expansion graph



Picture from Schaub et al. 2021

- Each interaction is transformed into a clique = all edges between pairs are present;
- HOIs actually disappeared !
- Too simplistic: For e.g, in co-authorship 1 paper with 3 authors \neq 3 different papers written by pairs of those authors.

Bipartite graph representation (two-modes network or star-expansion graph)



- No loss of information;
- But "higher-order" now translates into node degrees in one part;
- 2 two parts don't play symmetric roles: statistical models on bipartite graphs are not appropriate here

Picture from Schaub et al. 2021

Simple hypergraphs

Definition

A (simple) hypergraph $\mathcal{H} = (\mathcal{V}, \mathcal{E})$ is defined as a set of nodes $\mathcal{V} \neq \emptyset$ and a set of hyperedges \mathcal{E} . Each hyperedge is a non-empty collection of m distinct nodes ($2 \le m \le M$) taking part within an interaction.

- Hypergraphs naturally include the entity of graphs, by simply considering hyperedges of size m = 2;
- A hypergraph can contain a size-3 hyperedge [*a*, *b*, *c*] without any requirement on the existence of the size-2 hyperedges [*a*, *b*], [*a*, *c*], and [*b*, *c*].

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Clustering the nodes of a hypergraph I

What has been done up to now

Modularity-based approaches

- Different hypergraph modularity definitions: what kind of communities do they favour?
- Note that for computational reasons, these focus on *multisets-hypergraphs* where nodes may be repeated in a same hyperedge;
- > This is not always appropriate, e.g. co-authorship dataset;
- In the context of graphs, absence of self-loops and multiple edges are known to generate pbms in modularity approaches
- Spectral clustering has been generalized to hypergraphs but
 - it tends to favour groups of the same size;

Challenges

- Look for general clusters and not only communities
- None of these methods comes with a statistical criterion to select the number of groups *Q*

Clustering the nodes of a hypergraph II

Our proposal

- We focus on simple hypergraphs (instead of multisets-hypergraphs);
- We define a stochastic blockmodel to cluster the nodes of a hypergraph
 - We establish parameter identifiability results;
 - We propose a variational expectation-maximisation algorithm to infer clusters and parameters;
 - We propose an ICL criterion to select the number of clusters;
 - All these tools are implemented (in C++) in a efficient R **package** called HyperSBM.

Outline

2 Stochastic blockmodel for hypergraphs

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SBM formulation

- $\mathcal{H} = (\mathcal{V}, \mathcal{E})$, with $\mathcal{V} = \{1, \dots, n\}$ nodes and \mathcal{E} hyperedges;
- For each $2 \le m \le M$, let $\mathcal{V}^{(m)} = \{\{i_1, \dots, i_m\} : i_1, \dots, i_m \in \mathcal{V} \text{ and } i_1 \ne \dots \ne i_m\}$, set of unordered node tuples of size m;
- Observations: At each $\{i_1, \ldots, i_m\} \in \mathcal{V}^{(m)}$, we observe indicator variable $Y_{i_1,\ldots,i_m} = 1\{\{i_1,\ldots,i_m\} \in \mathcal{E}\};$
- Latent clusters: Z_1, \ldots, Z_n iid in $\{1, \ldots, Q\}$ with $\pi_q = \mathbb{P}(Z_i = q)$;
- Conditional independence assumption: $\{Y_{i_1,...,i_m}\}_{\{i_1,...,i_m\}\in\mathcal{V}^{(m)}}|\{Z_1,...,Z_n\}$ are independent with $Y_{i_1,...,i_m}|\{Z_1 = q_1,...,Z_m = q_m\} \sim \text{Bern}(B_{q_{i_1},...,q_{i_m}}^{(m)}).$

Parameter (generic) identifiability

Generic identifiability: a parameter θ almost surely (*w.r.t.* Lebesgue measure) uniquely defines the distribution \mathbb{P}_{θ} (up to label switching on the node groups).

Theorem

For any Q, the parameter $\theta = (\pi_q, B_{q_1,...,q_m}^{(m)})_{m,q,q_1,...,q_m}$ of the HSBM for (simple) hypergraphs over n nodes, is generically identifiable for large enough n.

Said differently, there is a finite set C of (non explicit) polynomial conditions on θ such that whenever $\theta \notin C$, the distribution \mathbb{P}_{θ} is uniquely defined by θ .

Inference through variational EM I

- Direct computation of the likelihood is not feasible for large *n*;
- EM algorithm neither feasible because latent variables are not independent conditional on observed ones;
- Variational approximation to EM algorithm: replace the intractable posterior distribution by the best approximation (w.r.t. Kullback-Leibler divergence) in a class of simpler (factorised) distributions:

$$\mathbb{Q}_{\tau}(Z_1,\ldots,Z_n)=\prod_{i=1}^n\mathbb{Q}_{\tau}(Z_i)=\prod_{i=1}^n\prod_{q=1}^Q\tau_{iq}^{Z_{iq}},$$

with the variational parameter $\tau_{iq} = \mathbb{Q}_{\tau}(Z_i = q) \in [0, 1]$ and $\sum_{q=1}^{Q} \tau_{iq} = 1$, for any i = 1, ..., n and q = 1, ..., Q.

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Inference through variational EM II

Evidence lower bound (ELBO)

$$egin{aligned} \mathcal{J}(heta, au) &= \mathbb{E}_{\mathbb{Q}_{ au}}[\log \mathbb{P}_{ heta}(oldsymbol{Y},oldsymbol{Z})] - \mathbb{E}_{\mathbb{Q}_{ au}}[\log \mathbb{Q}_{ au}(oldsymbol{Z})] \ &= \log \mathbb{P}_{ heta}(oldsymbol{Y}) - \mathsf{KL}(\mathbb{Q}_{ au}(oldsymbol{Z})||\mathbb{P}_{ heta}(oldsymbol{Z}|oldsymbol{Y})) \ &\leq \log \mathbb{P}_{ heta}(oldsymbol{Y}), \end{aligned}$$

with equality iff $\mathbb{Q}_{\tau}(\mathbf{Z})$ is the true posterior $\mathbb{P}_{\theta}(\mathbf{Z}|\mathbf{Y})$.

VEM maximises the lower bound $\mathcal{J}(\theta, \tau)$ (with respect to τ and θ) instead of the intractable log-likelihood log $\mathbb{P}_{\theta}(\mathbf{Y})$

VEM algorithm

• **VE-Step** maximizes $\mathcal{J}(\theta, \tau)$ with respect to τ :

$$\widehat{\tau}^{(t)} = \operatorname*{arg\,max}_{\tau} \, \mathcal{J}(\theta^{(t-1)}, \tau); \quad \text{s.t.} \ \sum_{q=1}^{Q} \tau_{iq} = 1 \qquad \forall i = 1, \dots, n.$$

This is equivalent to minimising the Kullback-Leibler divergence. In practice this step is obtained by a fixed-point algorithm.

• M-Step maximizes $\mathcal{J}(\theta, \tau)$ with respect to θ :

$$\widehat{\theta}^{(t)} = \operatorname*{arg\,max}_{\theta} \, \mathcal{J}(\theta, \tau^{(t-1)}), \quad \mathrm{s.t.} \ \sum_{q=1}^{Q} \pi_q = 1,$$

thus updating the value of the model parameters π_q and $B_{q_1,\ldots,q_m}^{(m)}$.

Model selection and generalizations

Integrated classification likelihood (ICL)

We select $\hat{q} = \arg \max_{q} ICL(q)$ where

$$\mathsf{ICL}(q) = \log \mathbb{P}_{\hat{\theta}}(\mathbf{Y}, \hat{\mathbf{Z}}) - \frac{1}{2}(q-1)\log n - \frac{1}{2}\sum_{m=2}^{M} \binom{q+m-1}{m}\log \binom{n}{m}.$$

Generalizations

- We have not considered self-loops (m = 1) but it's easy to do;
- Binary hyperedge variables could be replaced by counting hyperedges variables, replacing the Bernoulli distribution with, for e.g. (zero-inflated or deflated) Poisson law.

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Computational complexity - and considerations over the choice of M

- Focusing on *simple* hypergraphs has a high price: we need to explore all the ⁿ/_m tuples of nodes for all 2 ≤ m ≤ M;
- Our algorithm has a complexity of $O(n\binom{n}{M}Q^M)$, which is large;
- Current modularity approaches avoid this issue by working with multisets-hypergraphs, because there the summations over multisets of nodes $\sum_{i_1,...,i_m}$ factorize into *m* independent sums (no constraint that the nodes be different), and this further simplifies the expression of the modularity;
- Again, this is inappropriate on some datasets;
- As a consequence: we recommend to use a reasonable value of *M*: indeed *M* is not necessarily the largest observed hyperedge size (e.g. co-authorship dataset);

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Line clustering through hypergraphs I



Hypergraph construction

- Select 3 points at random and fit a line
- If residual distance less than a threshold, draw a hyperedge between those 3 points
- Globally set signal:noise hyperedge ratio = 2
- Repeat to obtain 100 3-uniform hypergraphs

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Line clustering through hypergraphs II

Data characteristics						
	Pts/line	Noisy pts	Total nb pts	mean nb of hyperedges		
2 lines	30	40	100	1070.84		
3 lines	30	60	150	587.7		

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Image: A matrix

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Comparison with modularity based methods I



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(a)

Comparison with modularity based methods II



Estimated number of groups

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Conclusions

- We propose a Stochastic Blockmodel for clustering the nodes of a (simple) hypergraph
- We establish (generic) identifiability of the parameters of the model
- Estimation and nodes clustering is performed through VEM algorithm
- ICL criterion is used to select the number of groups
- C++ code wrapped in a R package HyperSBM (https://github.com/LB1304/HyperSBM) and preprint on ArXiV https://arxiv.org/abs/2210.05983

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Remaining challenges

- understand the detectability limits for non-uniform hypergraphs ;
- computational issues: explore sparse hypergraphs modelings

Post-doc position on modelling sparse hypergraphs in Paris - deadline for application October, 15th.

Any questions ?

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Non equivalence between simple binary hypergraphs and bipartite graphs

Bipartite graphs space

Hypergraphs space

